

Comparaisons de génomes avec gènes dupliqués : étude théorique et algorithmes

Comparative genomics with duplicated genes: theoretical study and algorithms

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October 7th 2009

Outline

- 1 Genomes comparison
 - Overview
 - Genomes representation
 - Measures between genomes

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 - Exact approach
 - Heuristics and hybrid method

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 - Protocol
 - Visualization tool

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- 5 Conclusion

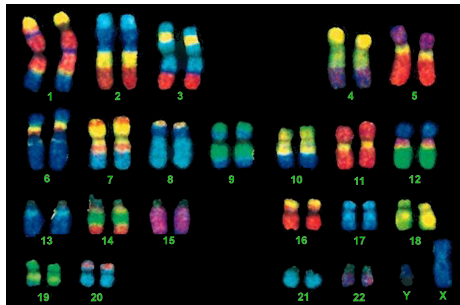
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Genomes and genes

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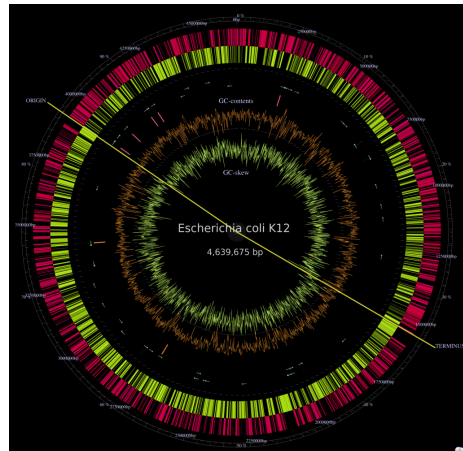
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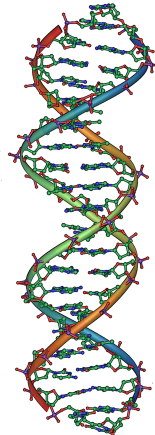
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Genomes and genes

Genome:

- Composed of one or several *chromosomes*
- Sequence(s) of *DNA*
- Hereditary information



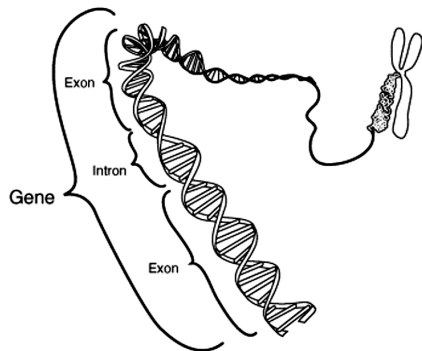
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Gene:

- Sequence of DNA
- Coding one or several *proteins*
- Gene orientation



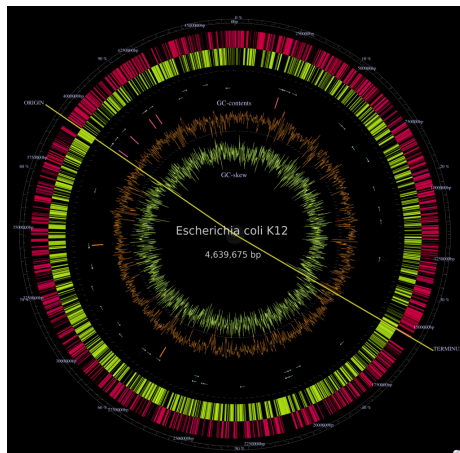
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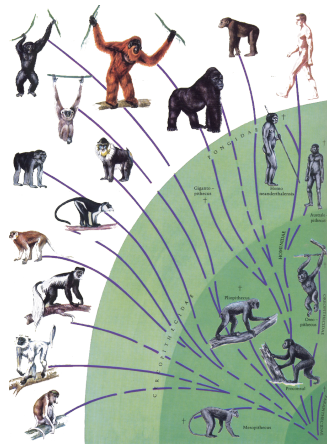
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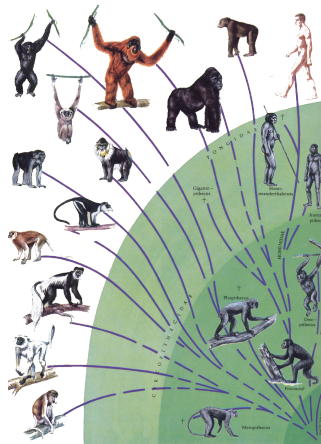
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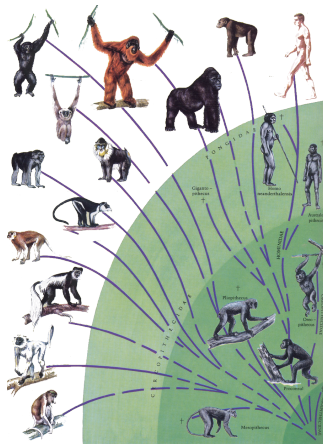
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Comparing genomes

Why?

- Phylogenetic trees construction
- Identification of highly conserved sequences
- Help genome annotation



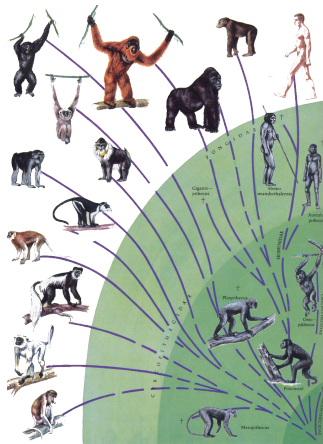
Comparing genomes

Why?

- Phylogenetic trees construction
- Identification of highly conserved sequences
- Help genome annotation

How?

- Genome modeled as a sequence of genes



Comparing two genomes : two different points of view

Comparison based on the evolution process

- Infer an evolution process from one genome to another
- Several operations can be considered:
 - ▶ inversion
 - ▶ translocation
 - ▶ duplication
 - ▶ ...
- Find a most parsimonious rearrangement scenario

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Comparison based on the structure of genomes

- Compare the structure (genes order) of the two genomes
- Compute a (dis)similarity measure between genomes
 - ▶ *number of breakpoints/adjacencies*
 - ▶ *number of common intervals*
 - ▶ *number of conserved intervals*
 - ▶ *Sum Adjacency Disruption*
 - ▶ ...

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Genomes representation

Representation and notations

- 1 Unichromosomal genome: sequence of *signed genes*

Example

- 1 $G_0 = +1 + 2 - 3 - 7 + 4 + 5 + 7 - 8 + 10 - 9 + 4 - 6 - 4$

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- 5 Let $\eta_{\mathbf{G}_0}$ be the number of genes in \mathbf{G}_0

Example

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- 2 $\Sigma = \{1, 2, 3 \dots 10\}$
- 3 $\mathbf{G}_0[4] = -7$
- 4 $\text{occ}(\mathbf{G}_0) = 3$
- 5 $\eta_{\mathbf{G}_0} = 13$

Measures between two genomes

- **Input:** Two genomes G_0 and G_1 with the same gene contents and without duplicates
 - **Output:** A (dis)-similarity measure between G_0 and G_1
-
- *number of breakpoints/adjacencies* [Watterson et al. 1982]
 - *number of common intervals* [Uno and Yagiura, 2000]
 - *number of conserved intervals* [Bergeron and Stoye, 2003]

Breakpoint and adjacency

Definition: **adjacency** and **breakpoint** [Watterson et al. 1982]

There exists an adjacency between genes $G_0[p]$ and $G_0[p + 1]$ iff $(G_0[p], G_0[p + 1])$ or $(-G_0[p + 1], -G_0[p])$ appears as a pair of consecutive genes in G_1 .

$$\begin{aligned}G_0 &= +1 + 2 + 3 + 4 + 5 \\G_1 &= +3 + 4 - 5 - 2 - 1\end{aligned}$$

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$$\begin{array}{l}
 \text{Adjacency} \\
 G_0 = \overbrace{+1 \ +2} + 3 + 4 + 5 \\
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 & & \text{Adjacency} & \text{Adjacency} & & \\
 & & \underbrace{\hspace{2em}} & \underbrace{\hspace{2em}} & & \\
 G_0 = & +1 & +2 & +3 & +4 & +5 \\
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Definition: **adjacency** and **breakpoint** [Watterson et al. 1982]

There exists a **breakpoint** between genes $G_0[p]$ and $G_0[p + 1]$ iff **neither** $(G_0[p], G_0[p + 1])$ **nor** $(-G_0[p + 1], -G_0[p])$ appears as a pair of consecutive genes in G_1 .

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 & & \overbrace{} & & \overbrace{} & & & & \\
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 G_1 = & -4 & -3 & & -5 & +1 & +2 & &
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$$\begin{array}{ccccccc}
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 \end{array}$$

Two measures:

- *Number of adjacencies*: similarity
- *Number of breakpoints*: dissimilarity

Common interval

Definition: **common interval** [Uno and Yagiura, 2000]

- A substring \mathbf{s}_0 of \mathbf{G}_0 is a *common interval* of $(\mathbf{G}_0, \mathbf{G}_1)$ if, in \mathbf{G}_1 , there is a substring \mathbf{s}_1 such that \mathbf{s}_1 is a permutation of \mathbf{s}_0 (without taking signs into account)

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- **Number of common intervals of $(\mathbf{G}_0, \mathbf{G}_1)$:**
Similarity measure between two genomes

Conserved interval

Definition: **conserved interval**

Proposed in [Bergeron and Stoye, 2003] for n permutations

- common interval
- same extremities OR reversed extremities

$$G_0 = +0 + 1 + 2 + 3 + 4 + 5$$

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- **Number of conserved intervals of (G_0, G_1) :**
Similarity measure between two genomes

And with duplicates?

- 1 Choose a one-to-one correspondence \mathcal{M} of genes (a matching)
- 2 Rename or remove genes according to \mathcal{M}
- 3 Compute the (dis)-similarity measure

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exemplar model (E)

[Sankoff, 99]

one occurrence for each
gene family in \mathcal{M}

$$G_0 = +0 +1 -2 -1 -3 +4$$

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$$G_0^E = +0 +1 -2 -3 +4$$

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maximum matching model (M)

[Tang & al, 03]

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$$G_0^M = +0 + 1' - 2 - 1'' - 3 + 4$$

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Intermediate model (I)

For each gene family,
at least one gene is kept in \mathcal{M}

Several possible matchings?

maximum matching model (M)

[Tang & al, 03]

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Measure between genomes with duplicates

Problem

- **Input:**
 - ▶ Two genomes G_0 and G_1
 - ▶ A model $X \in \{E, M, I\}$
- **Output:** Find a matching \mathcal{M} which satisfies the model X , and which optimizes the measure between G_0^X and G_1^X

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measure	problem
common interval	$ICOM_X$
conserved interval	$ICONS_X$
breakpoint	BD_X
adjacency	ADJ_X

Measure between genomes with duplicates

Problem

- **Input:**
 - ▶ Two genomes G_0 and G_1
 - ▶ A model $X \in \{E, M, I\}$
- **Question:** Are there G_0^X and G_1^X which satisfy the model X , and which imply no breakpoint ?

measure	problem	
common interval	$ICOM_X$	ZBD_X
conserved interval	$ICONS_X$	
breakpoint	BD_X	
adjacency	ADJ_X	

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What do we know?

	<i>exemplar model</i>	<i>maximum matching model</i>	<i>intermediate model</i>
$ICOM_x$ $ICONS_x$	NP-Complete [Chauve et al.] (instance (1, 2))		
BD_x	NP-Complete [Bryant] (instance (1, 2)) NP-Complete [Blin et al.] *		
ZBD_x	NP-Complete [Chen et al.] (instance (3, 3))	?	?

instance $(a, b) \Leftrightarrow occ(G_0) = a$ and $occ(G_1) = b$

* only one family contains several occurrences

Definition

α -approximation and PTAS

- Let P be an optimization problem
- Let I be an instance of P
- A polynomial algorithm A is an α -approximation iff
 - ▶ If P is a problem of minimization, then $A(I) \leq \alpha \cdot \mathit{optimal}(I)$
 - ▶ If P is a problem of maximization, then $A(I) \geq \frac{1}{\alpha} \cdot \mathit{optimal}(I)$

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- A polynomial algorithm B is a *Polynomial-Time Approximation Scheme* (PTAS) iff $\forall \epsilon > 0$
 - ▶ If P is a problem of minimization, then $B(I) \leq (1 + \epsilon) \cdot \mathbf{optimal}(I)$
 - ▶ If P is a problem of maximization, then $B(I) \geq \frac{1}{1+\epsilon} \cdot \mathbf{optimal}(I)$

APX-Hard Class

- If a problem P is **APX-Hard** then P does not admit a **PTAS**

New results

	<i>exemplar model</i>	<i>maximum matching model</i>	<i>intermediate model</i>
$ICOM_X$ $ICONS_X$	NP-Complete [Chauve et al.] (instance (1, 2)) APX-Hard (instance (1, 2)) *		
BD_X	NP-Complete [Bryant] (instance (1, 2))		NP-Complete [Blin et al.]
	APX-Hard (instance (1, 2)) *		
ZBD_X	NP-Complete [Chen et al.] (instance (3, 3)) (instance (2, k)) * [Blin et al.] (instance (2, 2))	polynomial *	$ZBD_I \equiv$ ZBD_E *
ADJ_X	$ADJ_E \simeq BD_E$ *	$ADJ_M \simeq BD_M$ *	$ADJ_I \neq BD_I$ *

* S. Angibaud, G. Fertin, I. Rusu, A. Thévenin et et S. Vialette
On the Approximability of Comparing Genomes with Duplicates

Journal of Graph Algorithms and Applications, Vol. 13(1), pages 19-53, 2009

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$A \simeq B$: An optimal solution for **A** is an optimal solution for **B**

$A \neq B$: An optimal solution for **A** is *not necessarily* an optimal solution for **B**

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⇒ Bad news : $ICOM_X$, $ICONS_X$ and BD_X do not admit a polynomial-time approximation scheme (PTAS)

New results

	<i>exemplar model</i>	<i>maximum matching model</i>	<i>intermediate model</i>
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⇒ Bad news : BD_E and BD_I do not admit any α -approximation, unless **P = NP**

New results

	<i>exemplar model</i>	<i>maximum matching model</i>	<i>intermediate model</i>
$ICOM_X$ $ICONS_X$	NP-Complete [Chauve et al.] (instance (1, 2)) APX-Hard (instance (1, 2))		
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ADJ_X	$ADJ_E \simeq BD_E$	$ADJ_M \simeq BD_M$	$ADJ_I \neq BD_I$

⇒ Good news : BD_M **could** admit an α -approximation

Outline

- 1 Genomes comparison
- 2 Theoretical complexity results
- 3 Algorithms
 - Exact approach
 - Pseudo boolean problem
 - Pseudo-boolean transformation for *ICOM_E*
 - Experimental results
 - Heuristics and hybrid method
 - IILCS_x
 - Hybrid method
 - Experimental results
- 4 MATCH&WATCH application

Exact algorithm

Problem

- **Input:**

- ▶ Two genomes G_0 and G_1
- ▶ A model $X \in \{E, M, I\}$

- **Output:** Find a matching \mathcal{M} which satisfies the model X , and which optimizes the measure between G_0^X and G_1^X

Idea: transformation into a pseudo boolean linear problem

Pseudo-boolean linear problem

Definition

- **Variables:** domain = $\{0, 1\}$
- **Constraints:** inequalities between weighted sum of variables
- **Objective function:** weighted sum of variables

Example

- **Variables:** $x \in \{0, 1\}, y \in \{0, 1\}, z \in \{0, 1\}$
- **Constraints:**
 - ▶ $x + 2 \cdot y \geq 2$
 - ▶ $z + y \leq 1$
- **Objective function:**
maximize $x + 2 \cdot y - z$

Pseudo-boolean linear problem

Definition

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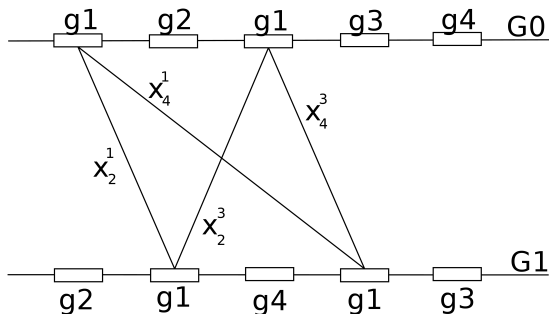
Example

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- **Constraints:**
 - ▶ $x + 2 \cdot y \geq 2$
 - ▶ $z + y \leq 1$
- **Objective function:**
maximize $x + 2 \cdot y - z$

⇒ Powerful solvers for this type of problem

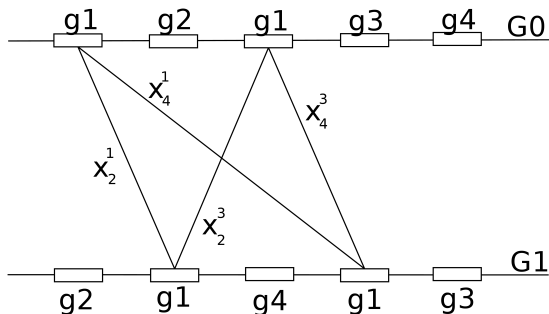
Transformation for $ICOM_E$: variables

- Variables x and I :



Transformation for $ICOM_E$: variables

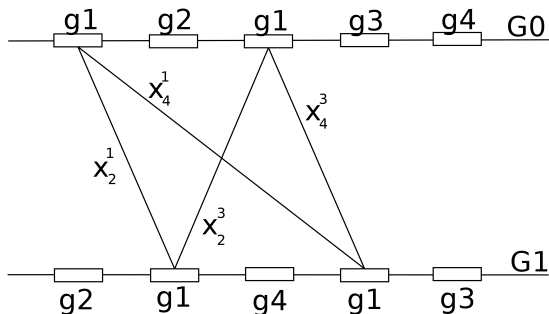
- Variables x and l :



x_b^a true \Leftrightarrow gene $G_0[a]$ and $G_1[b]$ are matched

Transformation for $ICOM_E$: variables

- Variables x and l :



$l_{k,l,m,n}$ true $\Leftrightarrow [k, l]$ in G_0 is a common interval of (G_0, G_1) , and $[m, n]$ in G_1 is a permutation of $[k, l]$

Transformation for $ICOM_E$: constraints

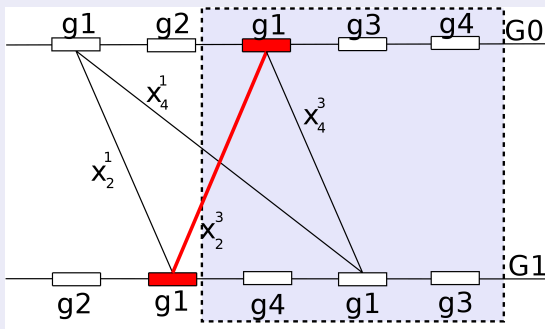
Exemplar model:

for each genome, only one occurrence of each gene family

$$C1: \forall f \in \mathcal{F}_{G_0} \cup \mathcal{F}_{G_1}, \quad \sum_{\substack{1 \leq a \leq \eta_{G_0} \\ G_0[a]=f}} \sum_{\substack{1 \leq b \leq \eta_{G_1} \\ G_1[b]=f}} x_b^a = 1$$

Transformation for $ICOM_E$: constraints

Validity of variables $l_{k,l,m,n}$



$$l_{k,l,m,n} + x_2^3 \leq 1$$

Transformation for $ICOM_E$

Objective function:

$$\text{Maximize } \sum_{k,l,m,n} I_{k,l,m,n}$$

Transformation for $ICOM_E$

Variables:

$$\mathcal{I} = \{I_{k,l,m,n} : 1 \leq k \leq \ell \leq \eta_{G_0} \wedge 1 \leq m \leq n \leq \eta_{G_1}\}$$

$$\mathcal{X} = \{x_b^a : 1 \leq a \leq \eta_{G_0} \wedge 1 \leq b \leq \eta_{G_1} \wedge G_0[a] = G_1[b]\}$$

Constraints:

$$(C.01) \quad \forall f \in \mathcal{F}_{G_0} \cup \mathcal{F}_{G_1}, \quad \sum_{\substack{1 \leq a \leq \eta_{G_0} \\ G_0[a]=f}} x_b^a = 1$$

$$(C.02) \quad \forall I_{k,l,m,n} \in \mathcal{I}, \quad \forall k < p < \ell, \quad \forall 1 \leq r < m, \quad G_0[p] = G_1[r], \quad I_{k,l,m,n} + x_r^p \leq 1$$

$$(C.03) \quad \forall I_{k,l,m,n} \in \mathcal{I}, \quad \forall k < p < \ell, \quad \forall n < r \leq \eta_{G_1}, \quad G_0[p] = G_1[r], \quad I_{k,l,m,n} + x_r^p \leq 1$$

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Transformation for $ICOM_E$

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Objective function:

$$\text{Maximize} \quad \sum_{k,l,m,n} I_{k,l,m,n}$$

Pseudo boolean transformation

Other problems ?

- **other models:** modify constraints C1
- **conserved intervals:** restriction on variables $I_{k,l,m,n}$
- **breakpoint and adjacency:** new variables and constraints

- **ICOM_X and ICONS_X**
S. Angibaud, G. Fertin, I. Rusu et S. Vialette.
A pseudo-boolean general framework for computing rearrangement distances between genomes with duplicates
Journal of Computational Biology, Vol. 14(4), pages 379-393. 2007

- **BD_X and ADJ_X**
S. Angibaud, G. Fertin, I. Rusu, A. Thévenin et S. Vialette.
Efficient Tools for Computing the Number of Breakpoints and the Number of Adjacencies between two Genomes with Duplicate Genes
Journal of Computational Biology, Vol. 15(8), pages 1093-1115. 2008

Experimental results

Dataset

- Twelve genomes of γ -*Proteobacteria* [Lerat et al. 2003]

Name	Genbank identifier	size
<i>Buchnera aphidicola</i> APS	NC_002528	564
<i>Escherichia coli</i> K12	NC_000913	4183
<i>Haemophilus influenzae</i> Rd	NC_000907	1709
<i>Pseudomonas aeruginosa</i> PA01	NC_002516	5540
<i>Pasteurella multocida</i> Pm70	NC_002663	2015
<i>Salmonella typhimurium</i> LT2	NC_003197	4203
<i>Wigglesworthia glossinidia</i> <i>brevipalpis</i>	NC_004344	653
<i>Xanthomonas axonopodis</i> pv. <i>citri</i> 306	NC_003919	4192
<i>Xanthomonas campestris</i>	NC_003902	4029
<i>Xylella fastidiosa</i> 9a5c	NC_002488	2680
<i>Yersinia pestis</i> CO_92	NC_003143	3599
<i>Yersinia pestis</i> KIM5 P12	NC_004088	3879

average: 3104

Experimental results

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- Twelve genomes of γ -*Proteobacteria* [Lerat et al. 2003]
- 66 possible pairs of genomes

Number of results:

	model		
	Exemplar	maximum matching	intermediate
ADJ_X	61/66	66/66	63/66
$ICOM_X$	21/66	40/66	21/66

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⇒ Efficient approach for ADJ_X

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⇒ Efficient approach for ADJ_X

⇒ Limit is attained for $ICOM_X$

⇒ Heuristics

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IILCS_M heuristic

- Based on ILCS_M heuristic [Tichy, 82]
- **Idea:** Match genes of a Longest Common Substring (LCS)

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IILCS_M heuristic

- 1 Compute the Longest Common Substring **S**

Example

$$\begin{array}{cccccccc}
 +1 & +2 & +3 & +4 & +5 & +6 & +7 & \\
 +6 & -7 & +4 & +5 & +1 & +6 & -3 & -2 & -1
 \end{array}$$

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- 2 Match all the genes of **S** accordingly

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IILCS_M heuristic

- 1 Compute the Longest Common Substring **S**
- 2 Match all the genes of **S** accordingly
- 3 **Remove genes that cannot be matched**
- 4 Iterate the process until saturation

Example

$$\begin{array}{ccccccc}
 +1 & +2 & +3 & +4 & +5 & +6 & +7 \\
 +6 & -7 & +4 & +5 & +6 & -3 & -2 & -1
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- Based on ILCS_M heuristic [Tichy, 82]
- **Idea:** Match genes of a Longest Common Substring (LCS)

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Example

$$\begin{array}{ccccccc}
 +1 & +2 & +3 & +4 & +5 & +6 & +7 \\
 +6 & -7 & +4 & +5 & +6 & -3 & -2 & -1
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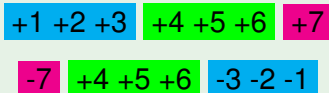
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- 5 Compute the measure

Example



Hybrid method

Algorithm $\text{HYB}_X(k)$

- **Idea:** Associate exact method and ILCS_X heuristic
- **Parameter k :** Bound on **LCS** size

1 Compute an **LCS** \mathbf{S} of (G_0, G_1)

2 **If** $|\mathbf{S}| \geq k$
Then

 Match all the genes of \mathbf{S}

 Remove genes that cannot be matched

 Return to 1

Else Apply the exact method: transformation into a pseudo-boolean linear problem

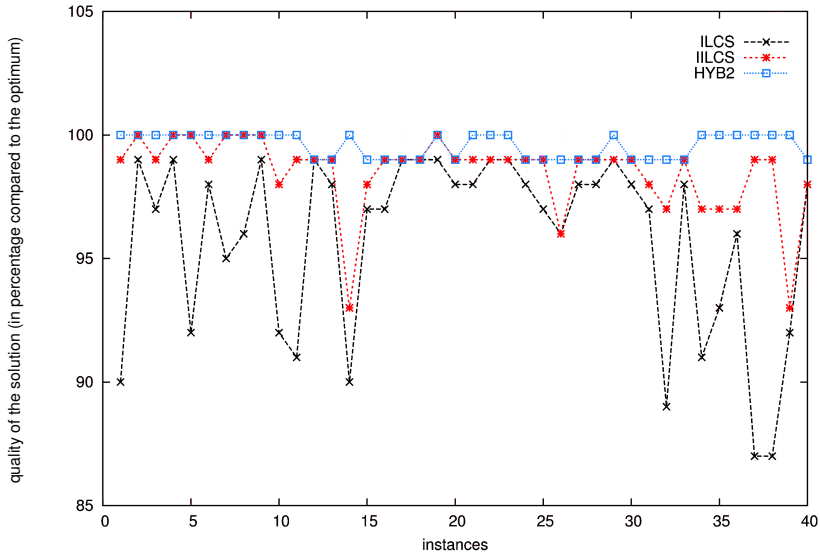
Experimental results

Dataset

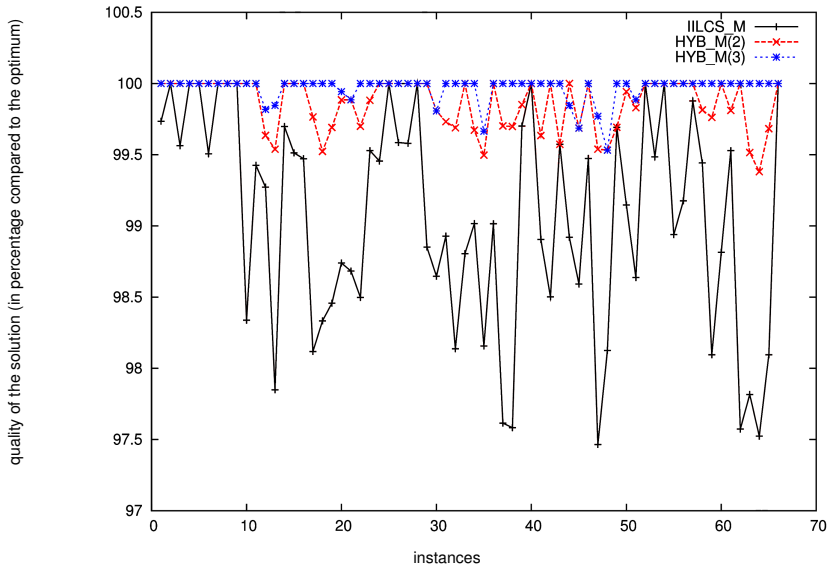
- Twelve genomes of γ -*Proteobacteria* [Lerat et al. 2003]
- 66 possible pairs of genomes

EXACT	model		
	Exemplar	maximum matching	intermediate
ADJ_X	61/66	66/66	63/66
$ICOM_X$	21/66	40/66	21/66

Experimental results: $ICOM_M$



Experimental results: ADJ_M



Outline

- 1 Genomes comparison
- 2 Theoretical complexity results
- 3 Algorithms
- 4 MATCH&WATCH application**
 - Protocol
 - Visualization tool
- 5 Conclusion

Goal

Problem

- **Input:** two circular genomes G_1 and G_2
- **Output:** List of common intervals between G_1 and G_2

Goal

- Compute common intervals
- Provide a tool to visualize and analyze results

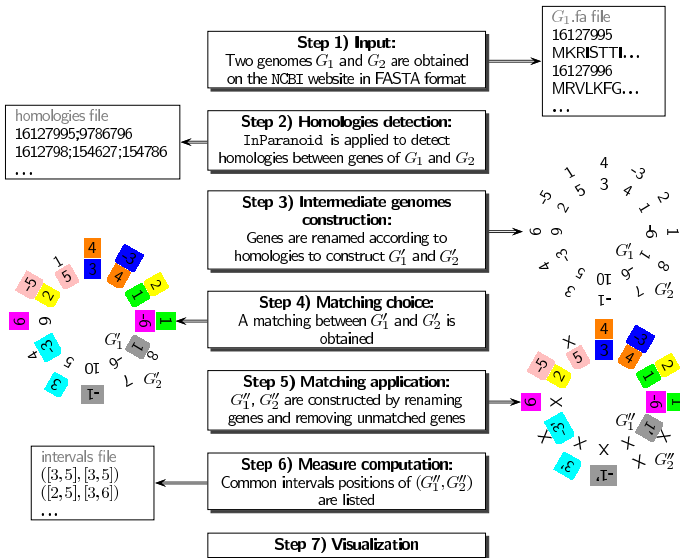
S. Angibaud, D. Éveillard, G. Fertin et I. Rusu

Comparing Bacterial Genomes by Searching Their Common Intervals

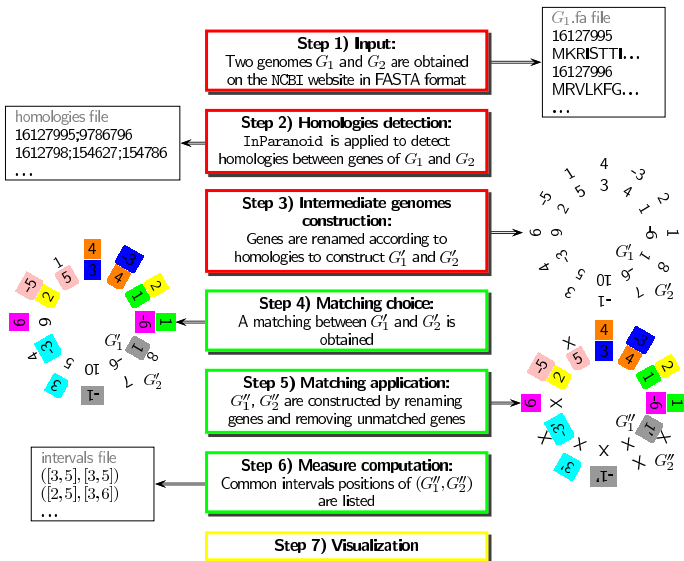
In Proc. 1st International Conference on Bioinformatics and Computational Biology

LNBI Vol. 5462, pages 102-113. 2009

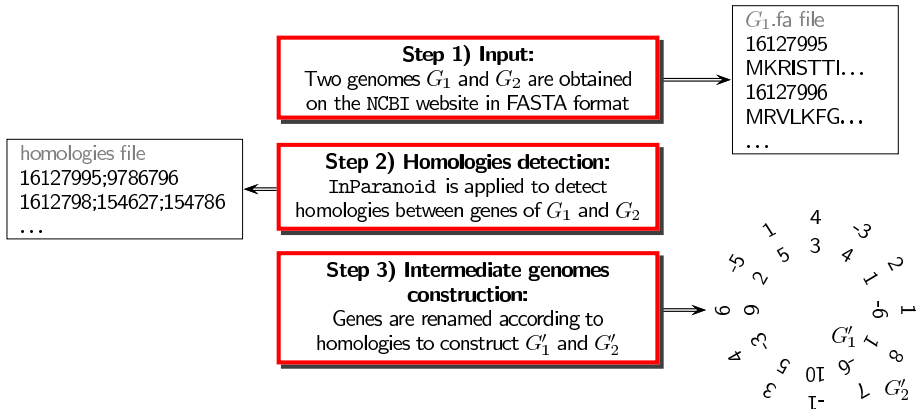
Protocol



Protocol



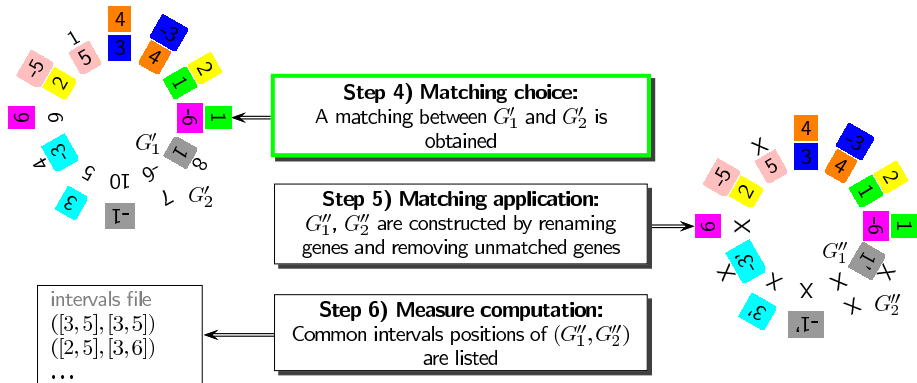
Homologies computation



Inparanoid [Storm et al. 2001]

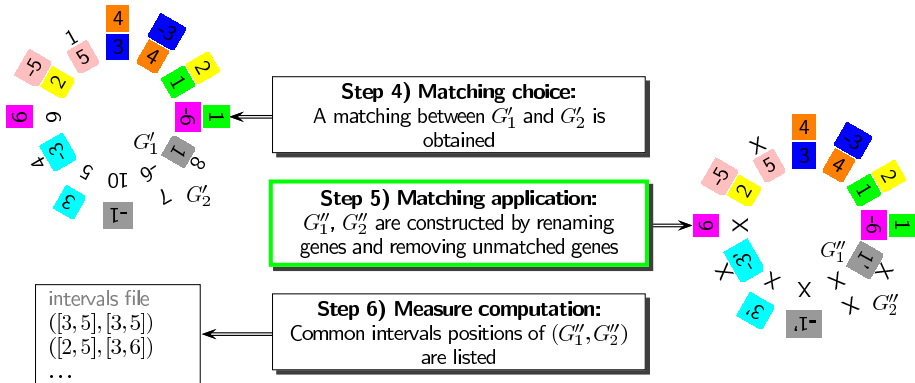
- Proposed in 2001 by Storm, Remm and Sonnhammer
- Compute clusters of homologous genes

Step 4: choose a matching

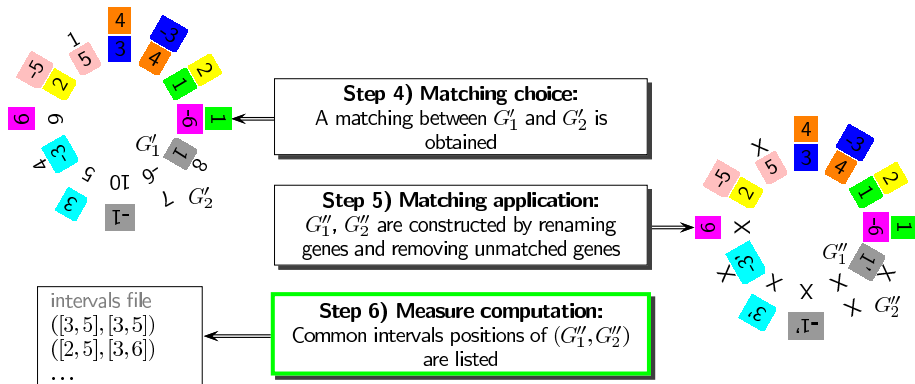


- Exact method: Pseudo boolean transformation
- IILCS_X heuristic
- Hybrid method

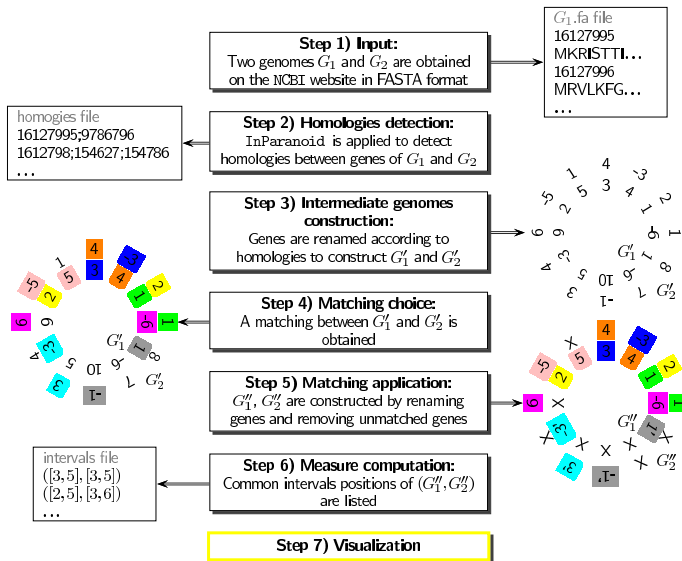
Step 5: Matching application



Step 6: common intervals computation



Seven steps



Outline

- 1 Genomes comparison
- 2 Theoretical complexity results
- 3 Algorithms
- 4 MATCH&WATCH application
- 5 Conclusion**

Contributions

- Better knowledge of problems
 - ▶ **APX**-Hardness of BD_X , $ICOM_X$ and $ICONS_X$
 - ▶ **NP**-Completeness of ZBD_E and ZBD_I
 - ▶ Polynomiality of ZBD_M

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- Three new algorithms
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 - Efficient approach for BD_X and ADJ_X
 - Limited for $ICOM_X$
 - ▶ IILCS $_X$ heuristic and Hybrid method
 - Promising results on a real dataset for each problem

Perspectives

- Work on MATCH&WATCH
 - ▶ First experimentation on six chromosomes of γ -Proteobacteria
 - ▶ Analyze in details the common intervals obtained
 - ▶ Add functionalities according to biologists

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 - ▶ α -approximation or PTAS for $ICOM_X$ on balanced genomes?

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- Partially ordered genomes

Acknowledgement

- *Directors*

- ▶ Irena Rusu
- ▶ Guillaume Fertin

- *Co-authors*

- ▶ Damien Éveillard (LINA, Université de Nantes)
- ▶ Annelyse Thévenin (LRI, Université Paris-Sud)
- ▶ Stéphane Vialette (IGM, Université Paris-Est Marne-la-Vallée)

Pictures

- http://www.mun.ca/biology/scarr/FISH_chromosomes_300dpi.jpg
- http://agaudi.files.wordpress.com/2008/09/dna_overview_es.png
- <http://joachimj.club.fr/imagesmada2004bis/PlanchePhylogenieidesprimates.jpg>
- <http://http://fr.wikipedia.org/wiki/Gene>
- <http://www.g-language.org/g3/>

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Thank you

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 - Pseudo boolean transformation for other problems
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Transformation for $ICOM_E$: objective function

Objective:

$$\text{maximize } \sum_{k,l,m,n} I_{k,l,m,n}$$

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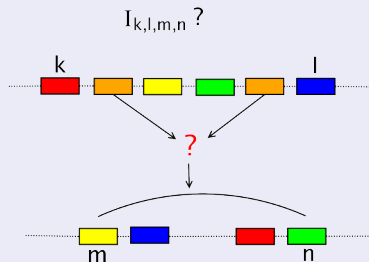
$$\text{maximize } \sum_{k,l,m,n} I_{k,l,m,n}$$

Improvements:

- Add rules to decrease the size of the instance

If all orange genes are located between the red and green one

We must have at least one orange gene to validate $I_{k,l,m,n}$



Transformation for $ICOM_E$: objective function

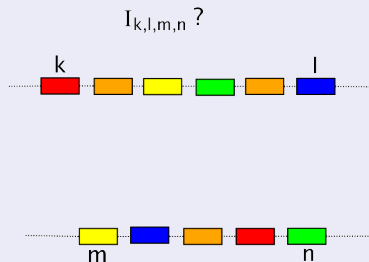
Objective:

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Improvements:

- Add rules to decrease the size of the instance

Else, we do not generate variable $I_{k,l,m,n}$



Other problems ?

Other models

- C1: (Exemplar model)

$$\forall f \in \mathcal{F}_{G_0} \cup \mathcal{F}_{G_1}, \quad \sum_{\substack{1 \leq a \leq \eta_{G_0} \\ G_0[a]=f}} \sum_{\substack{1 \leq b \leq \eta_{G_1} \\ G_1[b]=f}} x_b^a = 1$$

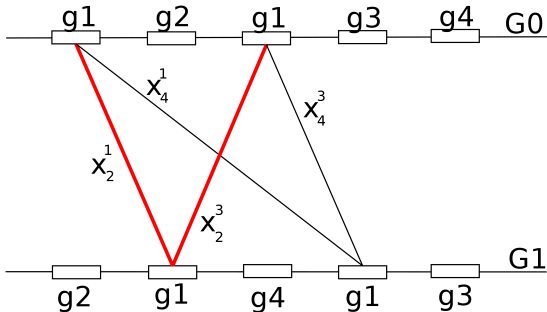
- C1': (Maximal matching model)

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- C1'': (Intermediate matching model)

$$\forall f \in \mathcal{F}_{G_0} \cup \mathcal{F}_{G_1}, \quad \sum_{\substack{1 \leq a \leq \eta_{G_0} \\ G_0[a]=f}} \sum_{\substack{1 \leq b \leq \eta_{G_1} \\ G_1[b]=f}} x_b^a \geq 1$$

Other models



- $\forall a = 1, 2, \dots, \eta_{G_0}, \quad \sum_{1 \leq b \leq \eta_{G_1}} x_b^a \leq 1$
 $G_0[a] = G_1[b]$
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Other measures

- $ICONS_X$:

Generate only variables $I_{k,l,m,n}$ such that

$$\left(\left(G_0[k] = G_1[m] \wedge G_0[l] = G_1[n] \right) \vee \right. \\ \left. \left(G_0[k] = -G_1[n] \wedge G_0[l] = -G_1[m] \right) \right)$$

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- BD_X and ADJ_X :

Other transformation

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LCS: Longest Common Substring [Tichy, 84]

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⇒ number of common intervals = 19

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⇒ number of common intervals = 20

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Heuristics: adaptation for other models

exemplar model

- For each gene family, we keep only the first occurrence in an LCS
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Heuristics: adaptation for other models

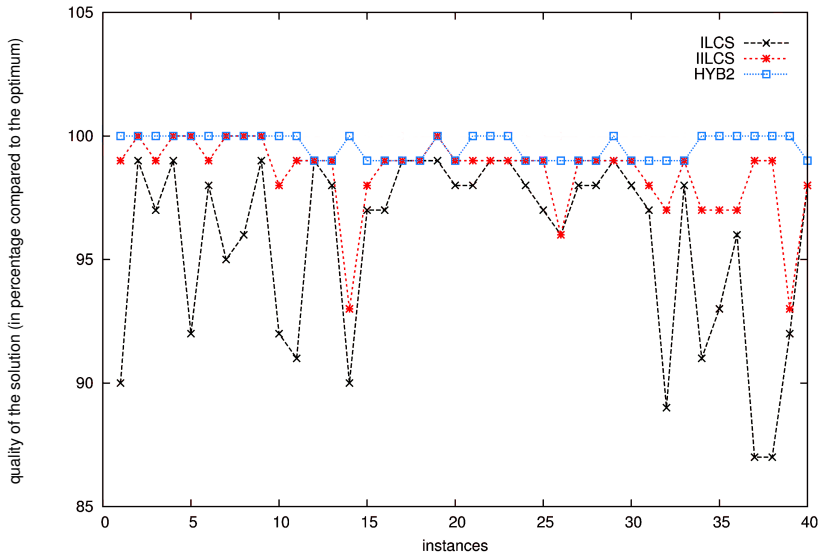
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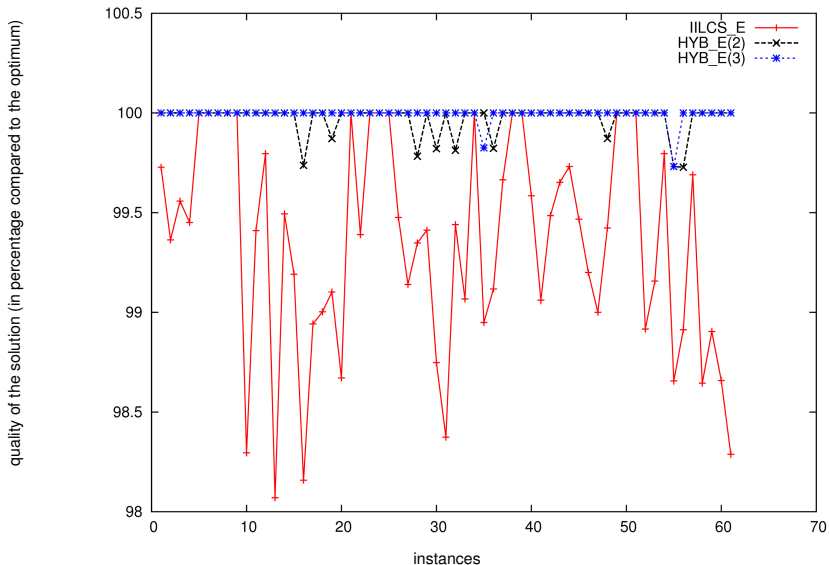
intermediate model

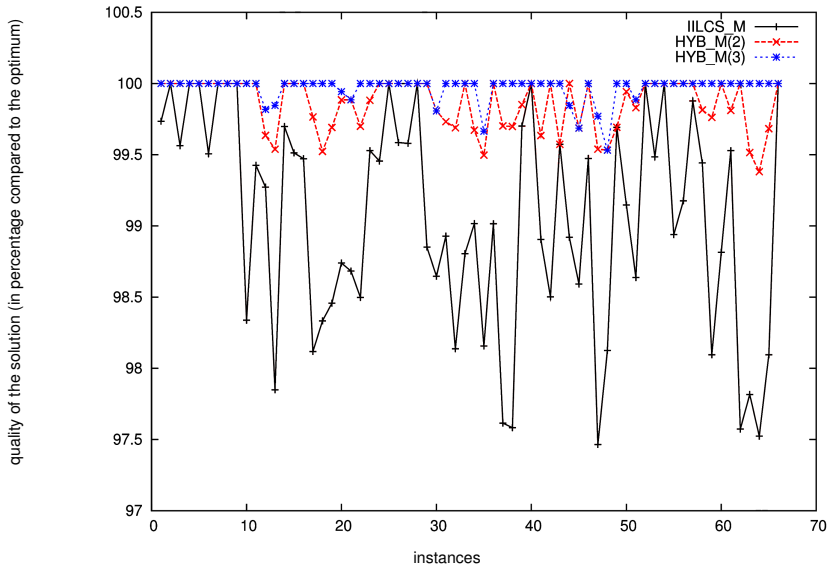
- We stop if, for each gene family, there exists at least one occurrence in the matching

Experimental results: $ICOM_M$

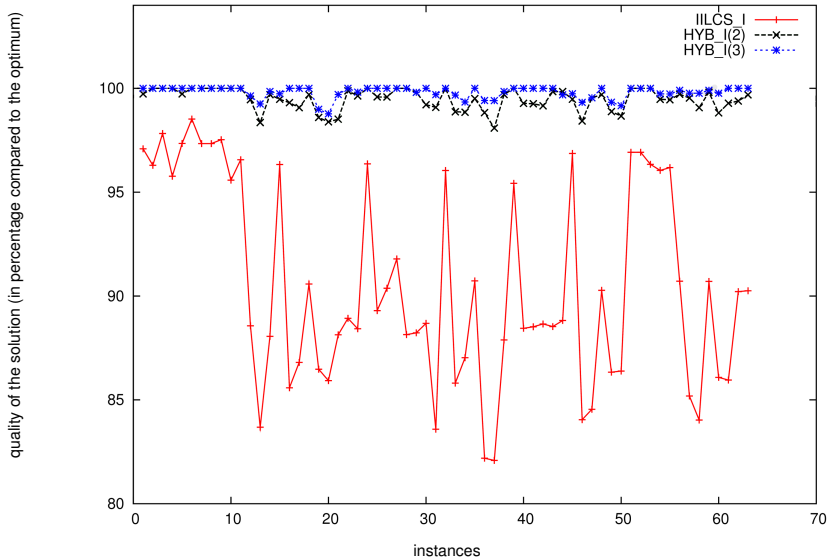


Experimental results: ADJ_E



Experimental results: ADJ_M 

Experimental results: ADJ_I



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Experience view

Experience Window

G1:

G2:

measure:

method:

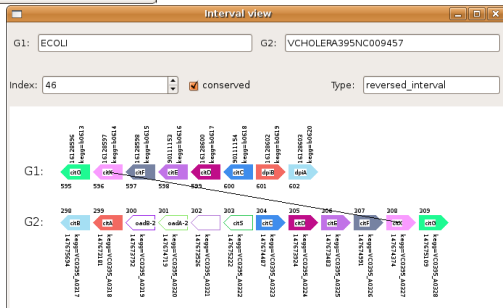
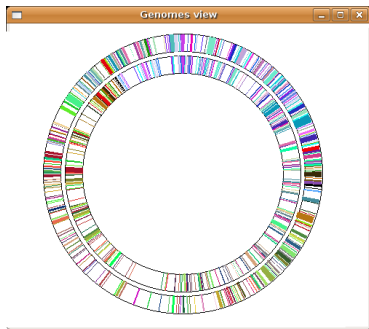
model:

Gene view

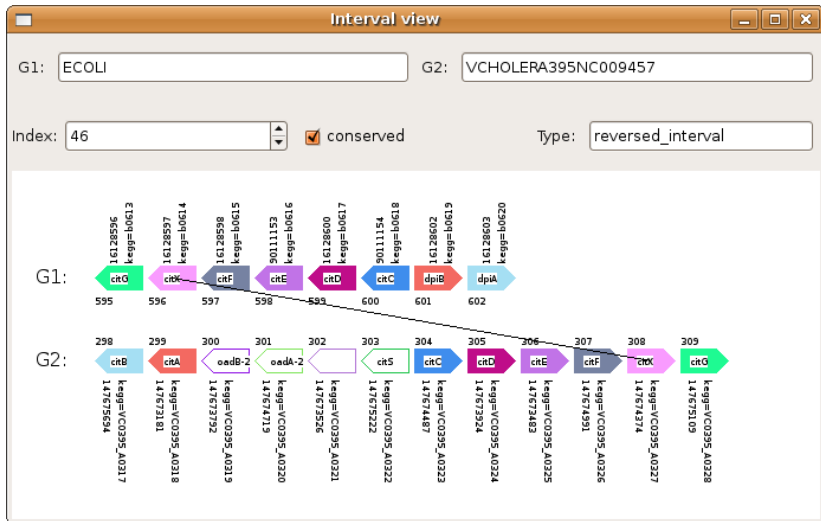
id: position:

genome:

Homologies		Pathways
<input type="text" value="ECOLI"/>	<input type="text" value="VCHOLERA1NC002505"/>	<input type="text" value="eco00670"/>
<input type="text" value="[47]16128042"/>	<input type="text" value="[434]15640467"/>	<input type="text" value="eco00790"/>



Visualization



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Common intervals filtering

- Lots of common intervals
 - Relevance of common intervals ?
- ⇒ Three filters to emphasize *the most interesting* common intervals

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Filters

- 1 **Maximal common intervals:**
Select only common intervals that are not contained in another one

Common intervals filtering

- Lots of common intervals
 - Relevance of common intervals ?
- ⇒ Three filters to emphasize *the most interesting* common intervals

Filters

1 **Maximal common intervals:**

Select only common intervals that are not contained in another one

2 **Annotated common intervals:**

Select maximal common intervals that contain some annotations in the *Ecocyc database*

Common intervals filtering

- Lots of common intervals
 - Relevance of common intervals ?
- ⇒ Three filters to emphasize *the most interesting* common intervals

Filters

1 Maximal common intervals:

Select only common intervals that are not contained in another one

2 Annotated common intervals:

Select maximal common intervals that contain some annotations in the *Ecocyc database*

3 Relevant common intervals:

Select annotated common intervals with good *p-value* (obtained by `GO-TermFinder`)

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Appendix

- 1 Appendix
 - Pseudo boolean transformation for other problems
 - $ILCS_X$ and $IILCS_X$
 - Visualization tool
 - Common intervals filtering
 - **First experimental results**

Experimental results

Input : six chromosomes of γ -Proteobacteria

NCBI identifiant	Name
NC_000913	<i>Escherichia coli</i> K12
NC_002505	<i>Vibrio cholerae</i> 01 biovar eltor str. N16961 chromosome I
NC_002506	<i>Vibrio cholerae</i> 01 biovar eltor str. N16961 chromosome II
NC_009456	<i>Vibrio cholerae</i> 0395 chromosome I
NC_009457	<i>Vibrio cholerae</i> 0395 chromosome II
NC_006840	<i>Vibrio fischeri</i> ES114 chromosome I
NC_006841	<i>Vibrio fischeri</i> ES114 chromosome II

Results: common intervals

genome G_2	genome size		method	computational time		common intervals	
	<i>E. coli</i>	G_2		Inparanoid (s)	matching (s)	number	maximal
NC002505	4243	2742	IILCS	1144	15	7418	274
NC002506	4243	1093	PSB	638	41	246	50
NC009456	4243	1133	PSB	651	46	264	55
NC009457	4243	2742	IILCS	1199	18	7204	278
NC006840	4243	2586	IILCS	1012	1	3865	255
NC006841	4243	1175	IILCS	715	1	203	62

Experimental results

